

Section 5

Lecturer: Larry Blume

TA: Abhi Ananth

Developed from Fikri Pitsuwan's notes.

1 Review

The *General Leontief Model (GLM)* extends the Simple Leontief Model (SLM) to capture the notion of multiple productive technology.

The main ingredients of the model are the $n \times m$ input matrix A and the $n \times m$ output matrix B , all with non-negative entries.

$$B = \begin{bmatrix} b_{11} & b_{12} & \dots & \dots & b_{1m} \\ \vdots & b_{22} & \dots & \dots & \vdots \\ \vdots & \vdots & \dots & \dots & \vdots \\ b_{n1} & \dots & \dots & \dots & b_{nm} \end{bmatrix} \begin{matrix} \text{Activity (m)} \\ \text{goods (n)} \end{matrix} \quad A = \begin{bmatrix} a_{11} & a_{12} & \dots & \dots & a_{1m} \\ \vdots & a_{22} & \dots & \dots & \vdots \\ \vdots & \vdots & \dots & \dots & \vdots \\ a_{n1} & \dots & \dots & \dots & a_{nm} \end{bmatrix}$$

Definition 1. (A, B) is productive if there exists an $x \geq 0$ such that $(B - A) \cdot x \gg 0$.

Assumptions maintained.

$n < m \quad |Z| = n \subseteq \{1, \dots, m\}$

- 1. (A, B) is productive
- 2. No joint production \rightarrow Any column/activity in B has at most 1 element > 0
- 3. Every good is produced \rightarrow Any row/good in B has at least 1 element > 0

Definition 2. A technology t is a set of n activities such that each good is produced by only one activity.

Lemma 1. If a GLM (A, B) is productive and produces $y \gg 0$, then there is a technology that produces y .

Consider the example:

$$B = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix} \quad A = \begin{bmatrix} 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 \end{bmatrix} \quad B - A = \begin{bmatrix} -1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 \end{bmatrix}$$

Take $y' = (1/2, 3/2)^T$, verify that both technologies $t = \{1, 2\}$ and $s = \{2, 3\}$ can produce y since

$$\begin{bmatrix} -1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 \end{bmatrix} \begin{bmatrix} 3 \\ 4 \\ 0 \end{bmatrix} = \begin{bmatrix} 1/2 \\ 3/2 \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} -1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \\ 3 \end{bmatrix} = \begin{bmatrix} 1/2 \\ 3/2 \end{bmatrix}$$

A natural question now is to ask which technology should we use? To compare technologies, we need some notion of efficiency of a technology. To that end, we introduce a finite resource, labor. For the GLM, the labor requirement vector is given by $a_0 = (a_{01}, \dots, a_{0m})$, where a_{0j} is the amount of labor needed to run

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$a_0 \cdot \vec{x} \rightarrow \text{scalar}$
 $\underline{a_0} \cdot \vec{x} \rightarrow \text{vector}$ $\begin{pmatrix} 0 \\ 1 \end{pmatrix}$ vs $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$.

→ one unit of activity j . Let L be the supply of labor. The set of *feasible net output (PPS)* for the GLM is $Y = \{y \in \mathbb{R}^n, y \geq 0 : y \leq (B - A)x, \underbrace{a_0 \cdot x}_{\text{labor}} \leq L, x \geq 0\}$.

The best way to produce y' is now the most efficient way to produce y' . Here, that will correspond to using the least amount of the finite resource, labor.

Example 1. Let $a_0 = (1, 1/2, 1/2)$, then to produce y' , technology t needs $a_0 \cdot (3, 4, 0)^T = 5$, while s needs $a_0 \cdot (0, 1, 3)^T = 2$. Thus, we should use technology s to produce y' .

Our analysis shows that some technologies are better than others, but what if we allow all activities to be used. Can we do any better? The answer is no.

For a given y , the following linear program gives the activity vector x that minimizes labor used.

$$\begin{aligned} \lambda(y) &= \min a_0 \cdot x \\ \text{s. t. } (B - A)x &\geq y \\ x &\geq 0 \end{aligned} \quad (\mathcal{L}(y))$$

For a given y and a technology t , the following linear program gives the activity vector using technology t that minimizes labor used.

$$\begin{aligned} \lambda^t(y) &= \min a_0 \cdot x \\ \text{s. t. } (B - A)x &\geq y \\ x &\geq 0 \\ x_i &= 0 \text{ if } i \notin t \end{aligned} \quad (\mathcal{L}^t(y))$$

First note that for a given y' , if x is an optimal solution to \mathcal{L} , then there is a basic optimal solution which is an optimal solution to \mathcal{L}^t for some t . That is, $\lambda(y') = \lambda^t(y')$ for some t . With the same reasoning, for a given y'' , there is a technology s such that $\lambda(y'') = \lambda^s(y'')$. There is no reason for s and t to be the same in our reasoning, but the next theorem says that they are.

Theorem 2 (Non-substitution). *If a GLM (A, B) satisfies the maintained assumptions, there is a technology t such that for all $y \geq 0$, $\lambda^t(y) = \lambda(y)$.*

Example 2. Let $y' = \begin{bmatrix} 1/2 \\ 3/2 \end{bmatrix}$ and $y'' = \begin{bmatrix} 1/2 \\ 1 \end{bmatrix}$. The problems are

$$\begin{aligned} \lambda(y') &= \min x_1 + \frac{1}{2}x_2 + \frac{1}{2}x_3 \\ \text{s. t. } -x_1 + x_2 &\geq 1 \\ x_1 + x_3 &\geq 3 \\ x &\geq 0 \end{aligned} \quad (\mathcal{L}(y'))$$

do x:
 $(B-A)x \geq y.$
 $x \geq 0.$

$$\begin{aligned} \lambda(y'') &= \min x_1 + \frac{1}{2}x_2 + \frac{1}{2}x_3 \\ \text{s. t. } -x_1 + x_2 &\geq 1 \\ x_1 + x_3 &\geq 4 \\ x &\geq 0 \end{aligned} \quad (\mathcal{L}(y''))$$

The solutions are $x' = (0, 1, 3)^T$ and $x'' = (0, 1, 4)^T$, respectively. We see that technology $s = \{2, 3\}$ gives $\lambda^s(y') = \lambda(y')$ and also that $\lambda^s(y'') = \lambda(y'')$.

$(0, 1)$ vs $(1, 0)$

$$\begin{aligned} \min & a_0 \cdot x \\ \text{st.} & (B - A)x \geq y \end{aligned}$$

Efficient production.

The non-substitution theorem says that technologies can be ranked by how efficient they are, where efficiency of a technology is measured by how much labor is used. Another interpretation of the non-substitution theorem is that when a GLM is productive, has one primary factor, and satisfies all our assumptions, restricting our usage to technology t with n activities does not alter the production possibility.

2 Problems

Problem 1 (Q7 of Linear Economic Models).

Problem 2. Consider a GLM given by

$$B = \begin{bmatrix} 1 & 1 \end{bmatrix} \quad A = \begin{bmatrix} 1/2 & 1/4 \end{bmatrix}$$

(and suppose that there are two primary factors $a_P = (1, 0)$ and $a_E = (0, 1/2)$ with labor supply $L_P = 4$ and $L_E = 2$, respectively. Is there a technology that does not alter the production possibility? Compare this to the case with one primary factor $a_0 = (1, 1/2)$ and $L = 6$.)

Tips:

- 1- PPS of all activities ?
- 2- PPS of all technologies ?

$(A, B) \rightarrow$ Productive ✓

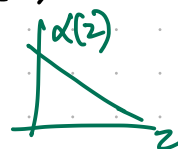
- ↳ No joint prod
- ↳ All goods produced.

7 Countries : A , B

Goods : $z \in [0, 1]$

Labor reqt : $(a(z), b(z))_{z \in [0, 1]}$ - only 1 input.

z is ordered so that $\alpha(z) = a(z)/b(z)$ strictly \downarrow in z .



Output prices : $(p(z))_{z \in [0, 1]}$

Input prices : w_a, w_b

Input endowments : L_A, L_B

a Convex support function of world PPS:

$$v(L_A, L_B) = \max_{x_a(\cdot), x_b(\cdot)} \int_0^1 p(z) \cdot [x_a(z) + x_b(z)] dz$$

$$\text{s.t.} \quad \int_0^1 a(z) x_a(z) dz \leq L_A \quad \text{----} (w_a)$$

$$\int_0^1 b(z) x_b(z) dz \leq L_B \quad \text{----} (w_b)$$

$$x_a(\cdot), x_b(\cdot) \geq 0$$

Complementary slackness:

$$w_a \left[\int_0^1 a(z) x_a(z) dz - L_A \right] = 0$$

$$w_b \left[\int_0^1 b(z) x_b(z) dz - L_B \right] = 0$$

$$\text{Dual} = \min_{\omega_a, \omega_b} \omega_a L_A + \omega_b L_B$$

$$\text{s.t. } \begin{cases} \omega_a \cdot a(z) \geq p(z) \quad \dots \quad (x_a(z)) \\ \omega_b \cdot b(z) \geq p(z) \quad \dots \quad (x_b(z)) \end{cases} \forall z$$

$\pi_a(z) \leq 0$ ←

$$\omega_a, \omega_b \geq 0$$

Profit maximization condition: $\forall z \in [0, 1]$,

$$x_a(z) \left[\underbrace{p(z) - \omega_a a(z)}_{=0} \right] = 0, \quad p(z) \leq \omega_a a(z)$$

$$x_b(z) \left[p(z) - \omega_b b(z) \right] = 0, \quad p(z) \leq \omega_b b(z)$$

[b] If $x_a(\tilde{z}) > 0$, $x_b(\tilde{z}) > 0$ for some $\tilde{z} \in [0, 1]$

then:

$$p(\tilde{z}) = \omega_a a(\tilde{z})$$

$$p(\tilde{z}) = \omega_b b(\tilde{z})$$

$$\Rightarrow \frac{a(\tilde{z})}{b(\tilde{z})} = \alpha(\tilde{z}) = \frac{\omega_b}{\omega_a}$$

Statement: $\exists z^*$: $\overset{A}{\leftarrow} B$ produced all $z < z^*$ & $\overset{A}{\leftarrow} B$ all $z > z^*$.

[c] Altered proposition: $\exists z^* \in [0, 1]$ such that

any $x_A^*(z) > 0$ has $z \geq z^*$ & any $x_B^*(z) > 0$ has $z \leq z^*$.
 not all $z < z^*$ has $x_B^*(z) > 0$

Pf: Key ingredient:

- i - Complementary slackness conditions
- ii - $\alpha(\cdot)$ is strictly decreasing

Part (b) above says that for any good produced in both countries, $\alpha(\tilde{z}) = \frac{w_b}{w_a}$.

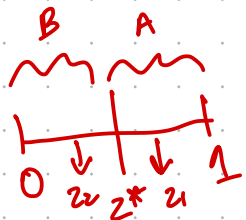
Since, $\alpha(\cdot)$ is strictly decreasing

$$\Rightarrow \exists \text{ at most } \underline{\text{one}} \tilde{z} \in [0, 1] : \alpha(\tilde{z}) = \frac{w_b}{w_a}$$

Towards a contradiction of proposition, suppose:

$$\exists z_1, z_2 \text{ such that } z_1 > z_2$$

Let A produces z_2 and B produces z_1 .



$$\begin{array}{l} \frac{w_a a(z_1) \geq p(z_1)}{w_a a(z_2) = p(z_2)} \quad \& \quad \frac{w_b b(z_1) = p(z_1)}{w_b b(z_2) \geq p(z_2)} \end{array}$$

$$\frac{a(z_1)}{a(z_2)} \geq \frac{p(z_1)}{p(z_2)} \quad \& \quad \frac{p(z_1)}{p(z_2)} \geq \frac{b(z_1)}{b(z_2)}$$

$$\Rightarrow \frac{a(z_1)}{b(z_1)} \geq \frac{a(z_2)}{b(z_2)} \equiv \alpha(z_1) \geq \alpha(z_2)$$

By strict decreasing: $\overset{z_1 > z_2}{\Rightarrow} \alpha(z_1) < \alpha(z_2) \Rightarrow \text{Contradiction}$

But suppose $x_A(z_1) > 0$, does this imply $x_A(z_2) > 0 \quad \forall z_2 > z_1$?

$$x_A(z_1) > 0 \Rightarrow p(z_1) = w_A a(z_1).$$

It is possible that:

$$p(z_2) < w_A a(z_2) \Rightarrow x_A(z_2) = 0$$

$$\frac{p(z_2)}{a(z_2)} < \frac{p(z_1)}{a(z_1)} !$$

$\rightarrow x_A(z_2) = 0$.
Not the case that all $z > z^*$ have $x_A(z) > 0$.

d) See b. Notice that z^* is the only good that both A and B can produce simultaneously.

This happens when: $0 \leq z^* = \frac{w_B}{w_A} \leq 1$

e) Identical consumers with Cobb-Douglas preferences

\Rightarrow Use aggregate representation.

Consumer problem:

$$\max_{C_A} U_A \left[(C_A(z))_{z \in [0,1]} \right].$$

$$\text{s.t.} \int_0^1 p(z) C_A(z) dz \leq \underbrace{w_A L_A}_{\text{selling labor in A.}}$$

$$s_i(z) = \frac{p(z) c_i(z)}{w_i L_i}$$

$U_A(\cdot)$ is Cobb Douglas

$$\Rightarrow \text{LNS} \Rightarrow \int_0^1 p(z) c_A(z) dz = w_A L_A$$

$$\therefore \int_0^1 p(z) \underbrace{\frac{c_A(z)}{w_A L_A}}_{\equiv s_A(z)} dz = 1$$

If $U_A = U_B$, given preferences are C-D:

$$\frac{p(z) c_A(z)}{w_A L_A} = \frac{p(z) c_B(z)}{w_B L_B}$$

$$\underline{s_A(z)} = \underline{s_B(z)} = s(z)$$

[f] Fraction spent on country A goods:

$$\theta(z) = \int_0^{z^*} s(z) dz \quad \tilde{\theta}(z^*) = \int_{z^*}^1 s(z) dz \quad (\because \tilde{\theta} = 1 - \theta \text{ in p.s.})$$

$$\text{Revenue in A} = \tilde{\theta}(z^*) [w_A L_A + w_B L_B]$$

$$\text{Labor expense in A} = w_A L_A$$

$$\text{Eqbm: } \tilde{\theta}(z^*) w_B L_B = w_A L_A [1 - \tilde{\theta}(z^*)]$$

$$\text{Revenue in B} = (1 - \tilde{\theta}(z^*)) [w_a L_a + w_b L_b]$$

$$\text{Labor expense in B} = w_b L_b$$

$$\text{Eqbm: } \tilde{\theta}(z^*) w_b L_b = (1 - \tilde{\theta}(z^*)) w_a L_a$$

$$\underbrace{\hspace{10em}}_X \quad \tilde{\theta} = \int_{z^*}^1 g(z) dz$$

$$\boxed{g} \quad \frac{w_b}{w_a} = \frac{L_a}{L_b} \frac{(1 - \tilde{\theta}(z^*))}{\tilde{\theta}(z^*)} \equiv \beta(z^*)$$

$$\boxed{h} \quad \beta(z^*) = \frac{w_b}{w_a} = \alpha(z^*)$$

$$\text{Diff. w.r.t. } L_a \quad \left(\frac{d\alpha(z^*)}{dz^*} = \frac{d\beta(z^*)}{dz^*} \right)$$

$$\alpha'(z^*) \cdot \frac{dz^*}{dL_a} = \frac{d}{dL_a} \left\{ \left(\frac{L_a}{L_b} \right) \left(\frac{1 - \tilde{\theta}(z^*)}{\tilde{\theta}(z^*)} \right) \right\}$$

$$= \frac{1}{L_b} \left[\frac{1 - \tilde{\theta}(z^*)}{\tilde{\theta}(z^*)} \right] + \frac{L_a}{L_b} \left[\frac{(-\tilde{\theta}(z^*) - (1 - \tilde{\theta}(z^*))) \frac{d\tilde{\theta}(z^*)}{dL_a}}{\tilde{\theta}(z^*)^2} \right]$$

$$= \frac{1}{L_b} \left[\frac{1 - \tilde{\theta}(z^*)}{\tilde{\theta}(z^*)} \right] - \frac{L_a}{L_b} \left(\frac{d\tilde{\theta}(z^*)/dz^* \cdot dz^*/dL_a}{\tilde{\theta}(z^*)^2} \right)$$

$$\frac{d\frac{u}{v}}{\frac{v}{v^2}} = \frac{v du - u dv}{v^2}$$

$$\tilde{\theta}(z^*) = \int_{z^*}^1 g(z) dz = \int_{z^*}^1 \underbrace{p(z)c(z)}_{w_A L_A + w_B L_B} dz \quad \geq 0$$

Using Leibniz Rule:

$$\frac{d\tilde{\theta}(z^*)}{dz^*} = -1 \cdot g(z^*)$$

$$\therefore \alpha'(z^*) \frac{dz^*}{dL_A} = \frac{1}{L_B} \left[\frac{1 - \tilde{\theta}(z^*)}{\tilde{\theta}(z^*)} \right] + \frac{L_A}{L_B} \left[\frac{g(z^*)}{\tilde{\theta}(z^*)^2} \right] \frac{dz^*}{dL_A}$$

$$\alpha'(z^*) L_B \tilde{\theta}(z^*)^2 \frac{dz^*}{dL_A} = \tilde{\theta}(z^*) (1 - \tilde{\theta}(z^*)) + L_A g(z^*) \frac{dz^*}{dL_A}$$

$$\frac{dz^*}{dL_A} = \frac{\tilde{\theta}(z^*) [1 - \tilde{\theta}(z^*)] \geq 0}{\underbrace{\alpha'(z^*) L_B \tilde{\theta}(z^*)^2}_{< 0} - \underbrace{L_A g(z^*)}_{> 0}}$$

$$\frac{dz^*}{dL_A} < 0$$

Intuition: $L_A \uparrow \Rightarrow w_A \downarrow \Rightarrow \frac{w_B}{w_A} \uparrow$,

$$\alpha(z^*) = \frac{w_B}{w_A} \text{ \& } \alpha'(\cdot) < 0$$

$$\Rightarrow z^* \downarrow$$

————— x —————

Q2

Start by defining the PPS using all technologies

$A_{n \times m}$

$n=1, m=2$

Each row of a_0 is a constraint.

$$Y = \{ y \geq 0 : (B-A)\vec{x} \geq y, \underline{a}_0 \vec{x} \leq \vec{L}, \vec{x} \geq 0 \}$$

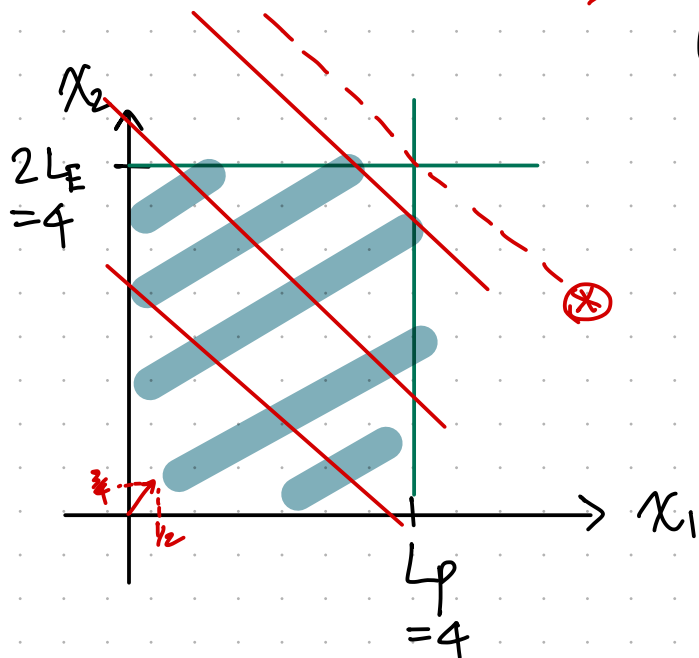
$B = [1, 1]$

$$\underline{a}_0 = \begin{bmatrix} \vec{a}_P \\ \vec{a}_E \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1/2 \end{bmatrix}; \quad \vec{L} = \begin{bmatrix} L_P \\ L_E \end{bmatrix} = \begin{bmatrix} 4 \\ 2 \end{bmatrix}$$

$$Y = \{ y \geq 0 : y \leq \left(\frac{x_1}{2} + \frac{3x_2}{4} \right), 0 \leq x_1 \leq 4$$

$$0 \leq \frac{x_2}{2} \leq 2 \}$$

$\underline{a}_0 : 1 \times m$



Set $x_1 = x_2 = 4$.

$$y = \frac{4}{2} + \frac{3 \times 4}{4} = 2 + 3 = \underline{\underline{5}}$$

Suppose we only use technology 1,

$$x_1 = 4, x_2 = 0 \Rightarrow y = \underline{\underline{2}}$$

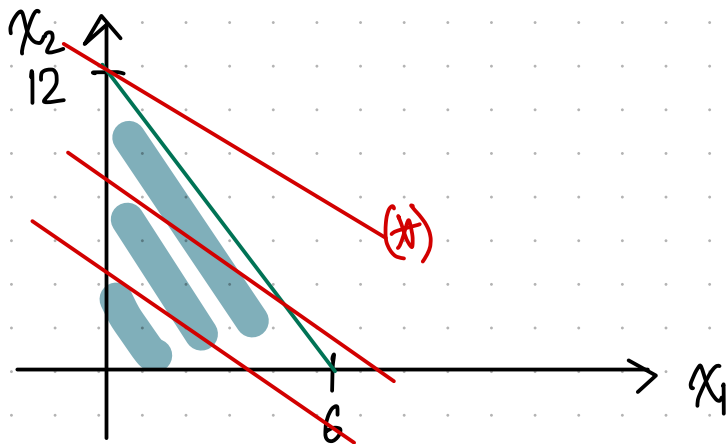
Suppose we only use technology 2,

$$x_1 = 0, x_2 = 4 \Rightarrow y = \underline{\underline{3}}$$

Thus, $\exists \tau \in \{\{1\}, \{2\}\}$ such that the PPS is unchanged

If $\vec{a}_0 = (1, 1/2)$; $L = 6$

$$Y = \left\{ y \geq 0 : y \leq \frac{x_1}{2} + \frac{3x_2}{4}, x_1 + \frac{x_2}{2} \leq 6, \vec{x} \geq 0 \right\}$$



$$x_1^* = 0; \quad x_2^* = \underline{\underline{12}}$$

$$y = \underline{\underline{9}}$$

$\tau = \{2\}$ produces same PPS.

